

On the order of recursive differentiability of finite binary quasigroups

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Abstract. The recursive derivatives of an algebraic operation are defined in [1], where they appear as control mappings of complete recursive codes. It is proved in [1], in particular, that the recursive derivatives of order up to r of a finite binary quasigroup (Q, \cdot) are quasigroup operations if and only if (Q, \cdot) defines a recursive MDS-code of length $r + 3$. The author of the present note gives an algebraic proof of an equivalent statement: a finite binary quasigroup (Q, \cdot) is recursively r -differentiable ($r \geq 0$) if and only if the system consisting of its recursive derivatives of order up to r and of the binary selectors, is orthogonal. This involves the fact that the maximum order of recursive differentiability of a finite binary quasigroup of order q does not exceed $q - 2$.

Mathematics subject classification: 20N05, 20N15, 11T71.

Keywords and phrases: quasigroup, recursive derivative, recursively differentiable quasigroup.

The notions of recursive derivative and recursively differentiable quasigroup have been introduced in [1], where the authors considered recursive MDS-codes (Maximum Distance Separable codes).

Let denote by $A^{(t)}$ the recursive derivative of order $t \geq 0$ of a binary groupoid (Q, A) , which is defined as follows:

$$A^{(0)} = A,$$

$$A^{(1)}(x, y) = A(y, A^{(0)}(x, y)),$$

$$A^{(t)}(x, y) = A(A^{(t-2)}(x, y), A^{(t-1)}(x, y)), \forall t \geq 2, \forall x, y \in Q.$$

A quasigroup (Q, A) is called *recursively r -differentiable* if the recursive derivatives $A^{(0)}, A^{(1)}, \dots, A^{(r)}$ are quasigroup operations ($r \geq 0$).

The notion of recursive derivative of a k -ary quasigroup (Q, A) , where $k \geq 2$, is defined in a similar way:

$$A^{(0)} = A,$$

$$A^{(t)}(x_1^k) = A(x_{t+1}, \dots, x_k, A^{(0)}(x_1^k), \dots, A^{(t-1)}(x_1^k)), \text{ if } 1 \leq t < k;$$

$$A^{(t)}(x_1^k) = A(A^{(t-k)}(x_1^k), \dots, A^{(k-1)}(x_1^k)), \text{ if } t \geq k, \forall x_1, \dots, x_k \in Q$$

(we denote by x_1^k the sequence x_1, x_2, \dots, x_k).

The length n of the codewords in a k -recursive code

$$C(n, A) = \{(x_1, \dots, x_k, A^{(0)}(x_1^k), \dots, A^{(n-k-1)}(x_1^k)) \mid x_1, \dots, x_k \in Q\}$$

given on an alphabet Q of q elements, where $A : Q^k \rightarrow Q$ is the defining k -ary operation, satisfies the condition $n \leq r + k + 1$, where r is the maximum order of the used recursive derivatives of (Q, A) . On the other hand, $C(n, A)$ is an MDS-code if and only if $d = n - k + 1$, where d is the minimum Hamming distance of this code. At present it is an open problem to determine all triplets (n, d, q) of natural numbers such that there exists an MDS-code C of length n , on an alphabet of q elements, with $|C| = q^k$ and with the minimum Hamming distance d , for each $k \geq 2$. This general question implies, in particular, the problem of determining the maximum order of recursive differentiability of finite k -ary quasigroups ($k \geq 2$).

It is known that there exist recursively 1-differentiable finite binary quasigroups of each order, excepting 1, 2, 6, and possibly 14, 18, 26 [1, 2]. Estimations of the maximum order r of recursive differentiability of finite n -quasigroups ($n \geq 2$) are given in [1, 3–6]. General properties of recursively differentiable binary quasigroups are studied in [5, 8].

The recursive differentiability of quasigroups is closely connected to the orthogonality of the recursive derivatives [1, 5, 8]. It is shown in [1] that a k -quasigroup defines an MDS-code of length n if and only if its first $n - k - 1$ recursive derivatives are strongly orthogonal. Hence the defining k -quasigroup operation of a recursive MDS-code of length n is recursively $(n - k - 1)$ -differentiable. On the other hand, it is known that a system of binary quasigroups is strongly orthogonal if and only if it is (simply) orthogonal [7]. It is proved in [1] that the recursive derivatives of order up to r of a finite binary quasigroup $(Q, *)$ are quasigroup operations if and only if $(Q, *)$ defines a recursive MDS-code of length $r + 3$.

In the present note we give an algebraic proof of the statement: a finite binary quasigroup $(Q, *)$ is recursively r -differentiable if and only if the system consisting of its recursive derivatives of order up to r is strongly orthogonal. This statement implies the fact that $r \leq q - 2$, where $q = |Q|$ and r is the maximum order of the recursive differentiability of the quasigroup Q .

Two binary operations A and B , defined on a set Q , are called orthogonal if the system of equations $A(x, y) = a, B(x, y) = b$ has a unique solution in Q , for every $a, b \in Q$. It follows from the previous definition that two binary operations A and B , defined on a set Q , are orthogonal if and only if the mapping

$$\sigma : Q \times Q \mapsto Q \times Q, \sigma(x, y) = (A(x, y), B(x, y))$$

is a bijection.

A system of binary operations $\{A_1, A_2, \dots, A_n\}$, $n \geq 2$, is said to be orthogonal if each two operations are orthogonal.

Denoting by F and E the binary selectors on a set Q : $F(x, y) = x$ and $E(x, y) = y$, $\forall x, y \in Q$, we get that a binary groupoid (Q, A) is a quasigroup if and only if A is orthogonal to each of two selectors.

Let (Q, A) be a binary quasigroup. It was observed by G. Belyavskaya [8] that $A^{(k)} = A\theta^k, \forall k \geq 1$, where $\theta = (E, A)$. An analogous representation for the recursive derivatives of k -ary operations ($k \geq 2$) was given in [5].

Theorem 1. *A finite binary quasigroup (Q, A) is recursively n -differentiable if and only if the system $\{F, E, A, A^{(1)}, \dots, A^{(n)}\}$ is orthogonal.*

Proof. Let (Q, A) be a recursively n -differentiable finite binary quasigroup. Then the recursive derivatives $A^{(1)}, \dots, A^{(n)}$ are quasigroup operations, so each recursive derivative $A^{(k)}$ of the system is orthogonal to the selectors F and E .

Now, let k and s be two distinct numbers between 0 and n : $0 \leq k < s \leq n$. As

$$(A^{(k)}, A^{(s)}) = (A\theta^k, A\theta^s) = (A, A^{(s-k)})\theta^k,$$

where $\theta = (E, A)$ is a bijection, we get that $A^{(k)}$ and $A^{(s)}$ are orthogonal if and only if A and $A^{(s-k)}$ are orthogonal, i.e. if and only if A and $A^{(m)}$ are orthogonal, for every $m = 1, 2, \dots, n$. On the other hand,

$$A^{(m)}(x, y) = A^{(m-1)}(E, A)(x, y) = A^{(m-1)}(y, A(x, y)),$$

hence the system of equations

$$\begin{cases} A(x, y) = a, \\ A^{(m)}(x, y) = b, \end{cases}$$

is equivalent to

$$\begin{cases} A(x, y) = a, \\ A^{(m-1)}(y, a) = b, \end{cases}$$

which has a unique solution as A and $A^{(m-1)}$ are quasigroup operations. Therefore the system $\{F, E, A, A^{(1)}, \dots, A^{(n)}\}$ is orthogonal.

Conversely, if the system $\{F, E, A, A^{(1)}, \dots, A^{(n)}\}$ is orthogonal, then each of the recursive derivatives $A, A^{(1)}, \dots, A^{(n)}$ is orthogonal to the selectors F and E , hence the recursive derivatives of order up to n are quasigroup operations, i.e. (Q, A) is recursively n -orthogonal. □

Corollary 1. *The maximum order r of recursive differentiability of a finite binary quasigroup of order q does not exceed $q - 2$.*

Proof. The proof follows from the fact that there exist at most $q - 1$ pairwise orthogonal latin squares of order q , which implies that the maximum order r of recursive differentiability satisfies the inequality $r + 1 \leq q - 1$, hence $r \leq q - 2$. □

It is shown in [1] that there exist recursively $(q - 2)$ -differentiable finite binary quasigroups of every primary order $q \geq 3$. However, it is an open problem to find the maximum order of recursive differentiability of finite k -ary quasigroups of order q , for $k \geq 2$ and an arbitrary non-primary q .

Acknowledgment. This work is partially supported by National Agency for Research and Development of the Republic of Moldova, project 20.80009.5007.25.

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Received July 21, 2022